6.033 Computer System Engineering

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L11: Link and Network layer

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Some slides are from lectures by Nick Mckeown, Ion Stoica, Frans Kaashoek, Hari Balakrishnan, Sam Madden, and Robert Morris

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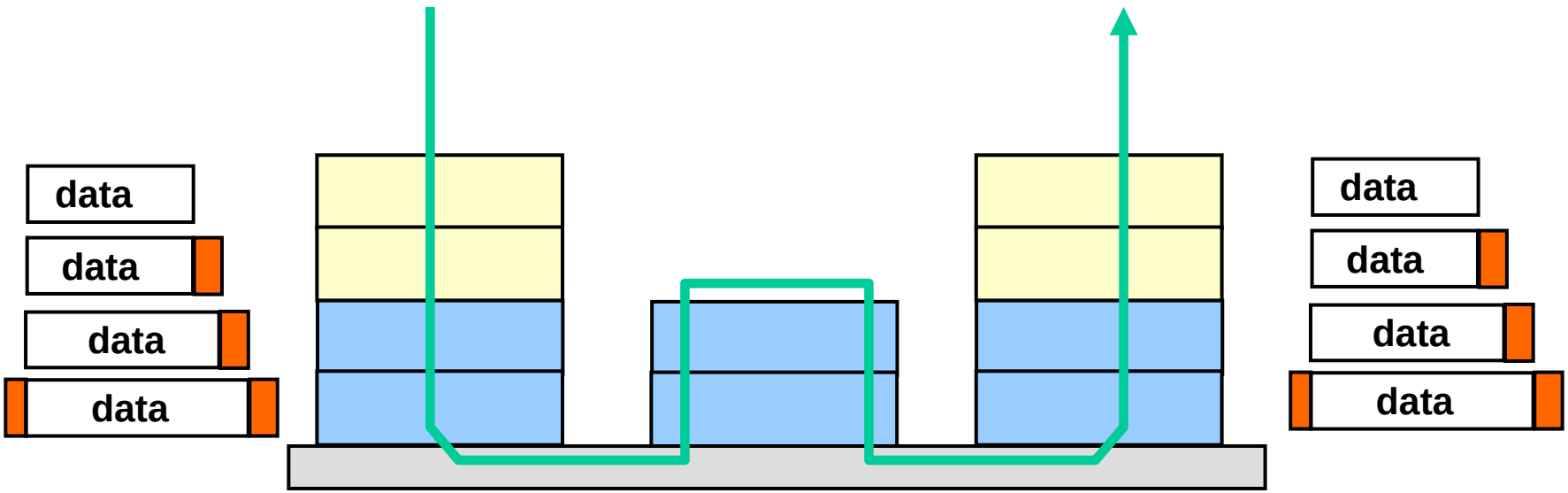
Last lecture: layering of protocols

- Each layer adds/strips off its own header
- Each layer may split up higher-level data
- Each layer multiplexes multiple higher layers
- Each layer is (mostly) transparent to higher layers



Link Layer

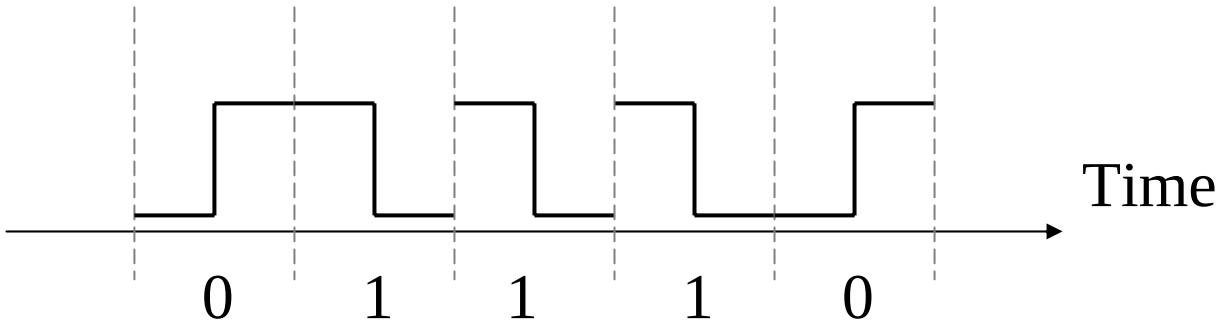
Problem:

Deliver data from one end of the link to the other

Need to address:

- Bits→ Analog → Bits
- Framing
- Errors
- Medium Access Control (The Ethernet Paper)

Manchester encoding



- Each bit is a transition
- Allows the receiver to sync to the sender's clock

Framing

- Receiver needs to detect the beginning and the end of a frame
- Use special bit-pattern to separate frames
  - E.g., pattern could be 1111111 (7 ones)
- *Bit stuffing* is used to ensure that a special pattern does not occur in the data
  - If pattern is 1111111 → Whenever the sender sees a sequence of 6 ones in the data, it inserts a zero (reverse this operation at receiver)

Error Handling

- Detection:
  - Use error detection codes, which add some redundancy to allow detecting errors
- When errors are detected
  - Correction:
    - Some codes allow for correction
  - Retransmition:
    - Can have the link layer retransmit the frame (rare)
  - Discard:
    - Most link layers just discard the frame and rely on higher layers to retransmit

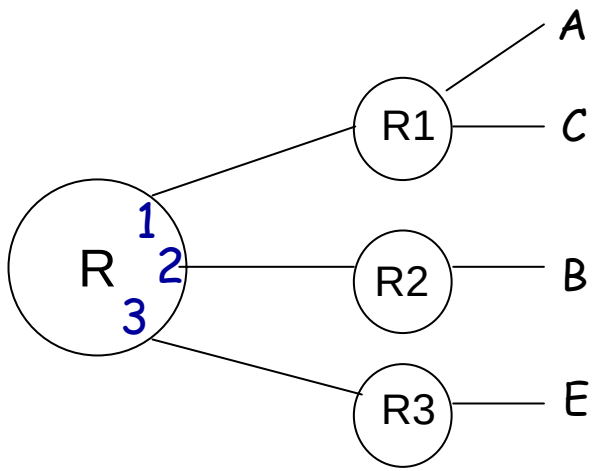
Network Layer:

finds a path to the destination and forwards packets along that path

- Difference between routing and forwarding
  - Routing is finding the path
  - Forwarding is the action of sending the packet to the next-hop toward its destination

Forwarding

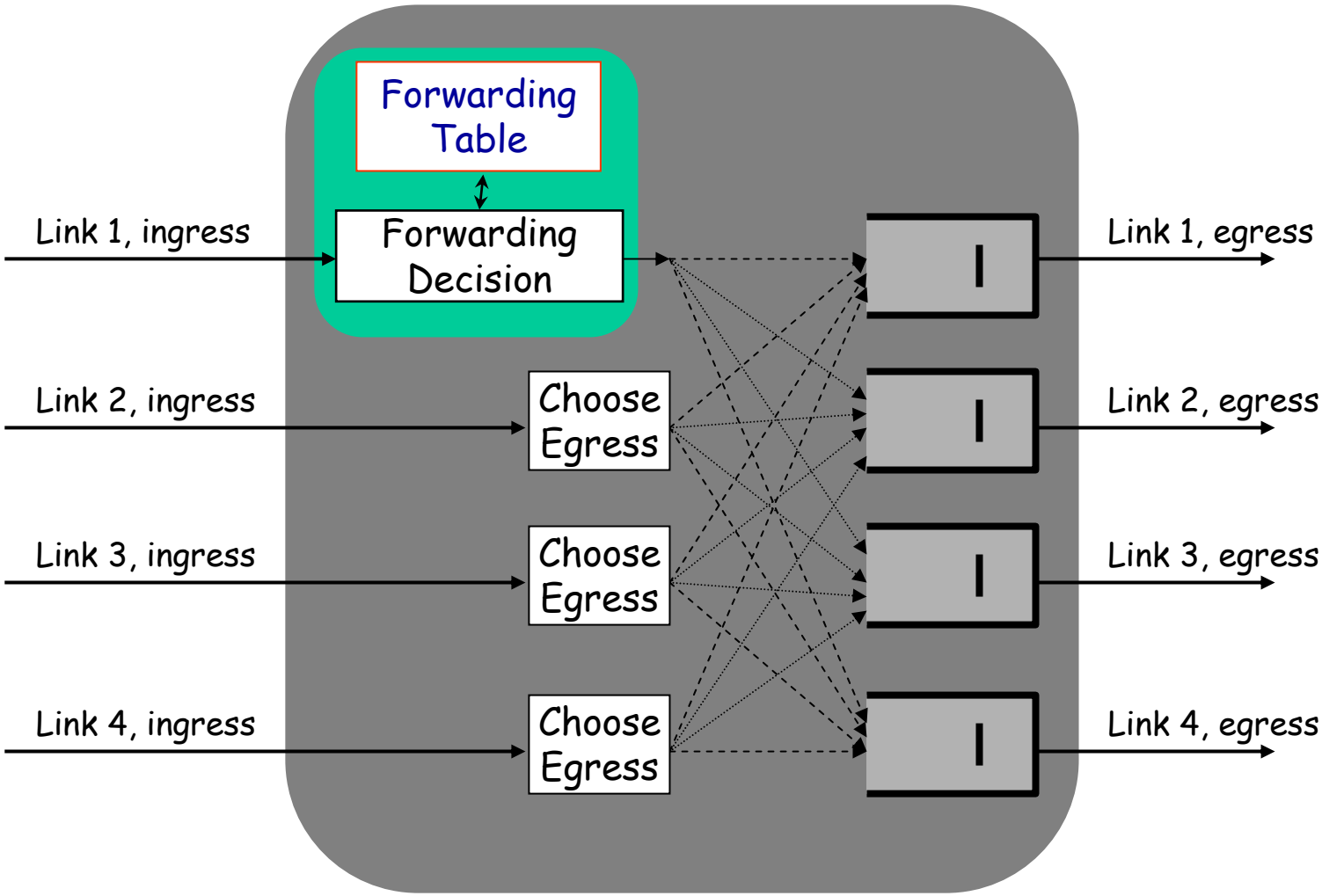
- Each router has a forwarding table
- Forwarding tables are created by a routing protocol



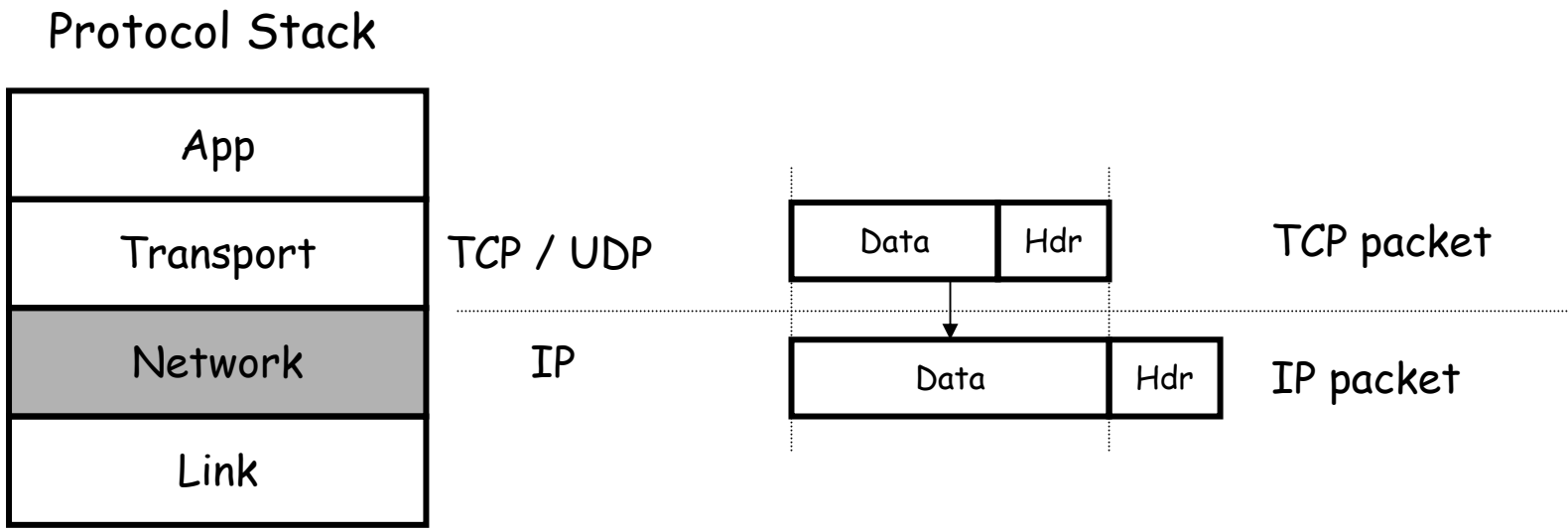
Forwarding table at R

Dst. Addr	Link
A	1
B	2
C	1
E	3

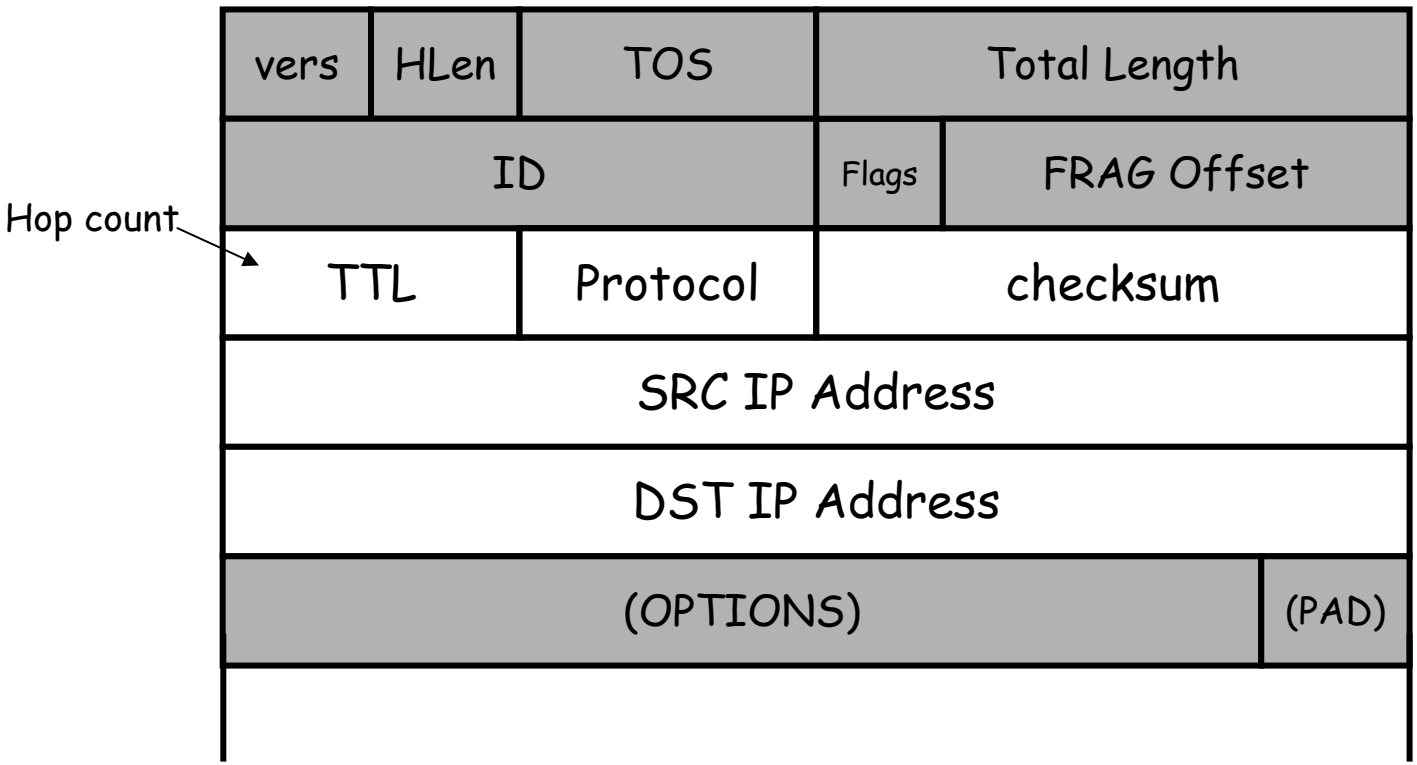
Inside a router



The Internet Protocol (IP)



The IP Header

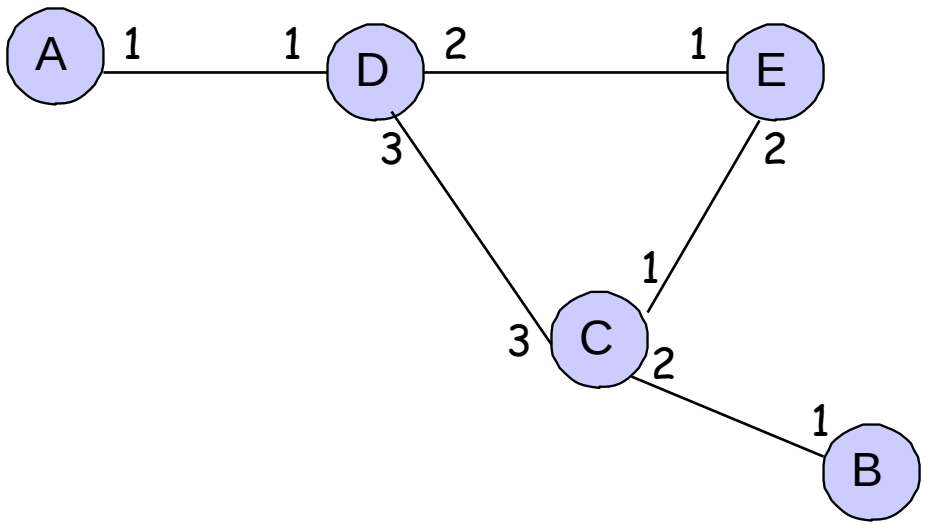


Forwarding an IP Packet

- Lookup packet's DST in forwarding table
  - If known, find the corresponding outgoing link
  - If unknown, drop packet
- Decrement TTL and drop packet if TTL is zero; update header Checksum
- Forward packet to outgoing port
- Transmit packet onto link

The Routing Problem:

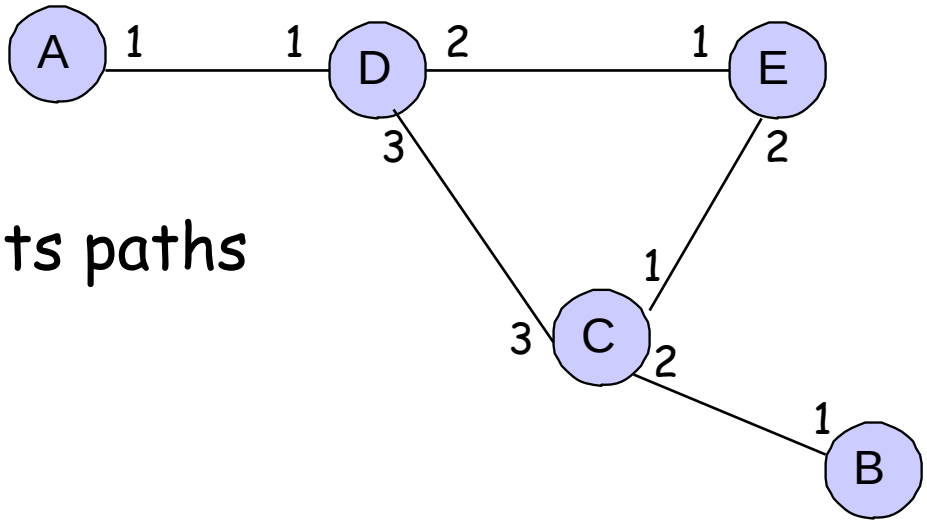
- Generate forwarding tables



Goals: No loops, short paths, etc.

Path Vector Routing Protocol

- Initialization
  - Each node knows the path to itself

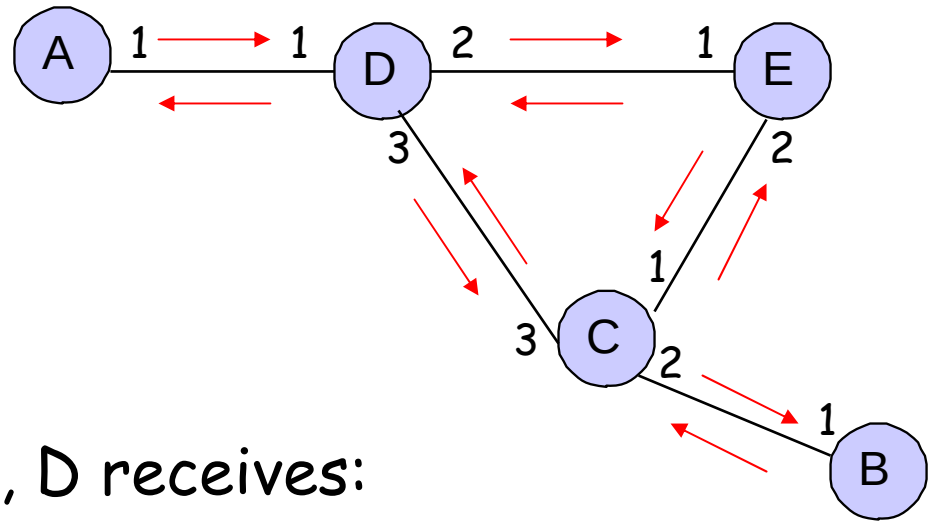


For example, D initializes its paths

DST	Link	Path
D	End layer	null

Path Vector

- Step 1: Advertisement
  - Each node tells its neighbors its path to each node in the graph



For example, D receives:

From A:

To	Path
A	null

From C:

To	Path
C	null

From E:

To	Path
E	null

Path Vector

- Step 2: Update Route Info
  - Each node use the advertisements to update its paths

D received: From A:

To	Path
A	null

From C:

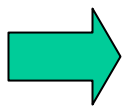
To	Path
C	null

From E:

To	Path
E	null

D updates its paths:

DST	Link	Path
D	End layer	null



DST	Link	Path
D	End layer	null
A	1	<A>
C	3	<C>
E	2	<E>

Note: At the end of first round, each node has learned all one-hop paths

Path Vector

- Periodically repeat Steps 1 & 2

In round 2, D receives:

From A:

To	Path
A	null
D	<D>

From C:

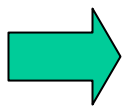
To	Path
C	null
D	<D>
E	<E>
B	<B>

From E:

To	Path
E	null
D	<D>
C	<C>

D updates its paths:

DST	Link	Path
D	End layer	null
A	1	<A>
C	3	<C>
E	2	<E>



DST	Link	Path
D	End layer	null
A	1	<A>
C	3	<C>
E	2	<E>
B	3	<C, B>

Note: At the end of round 2, each node has learned all two-hop paths

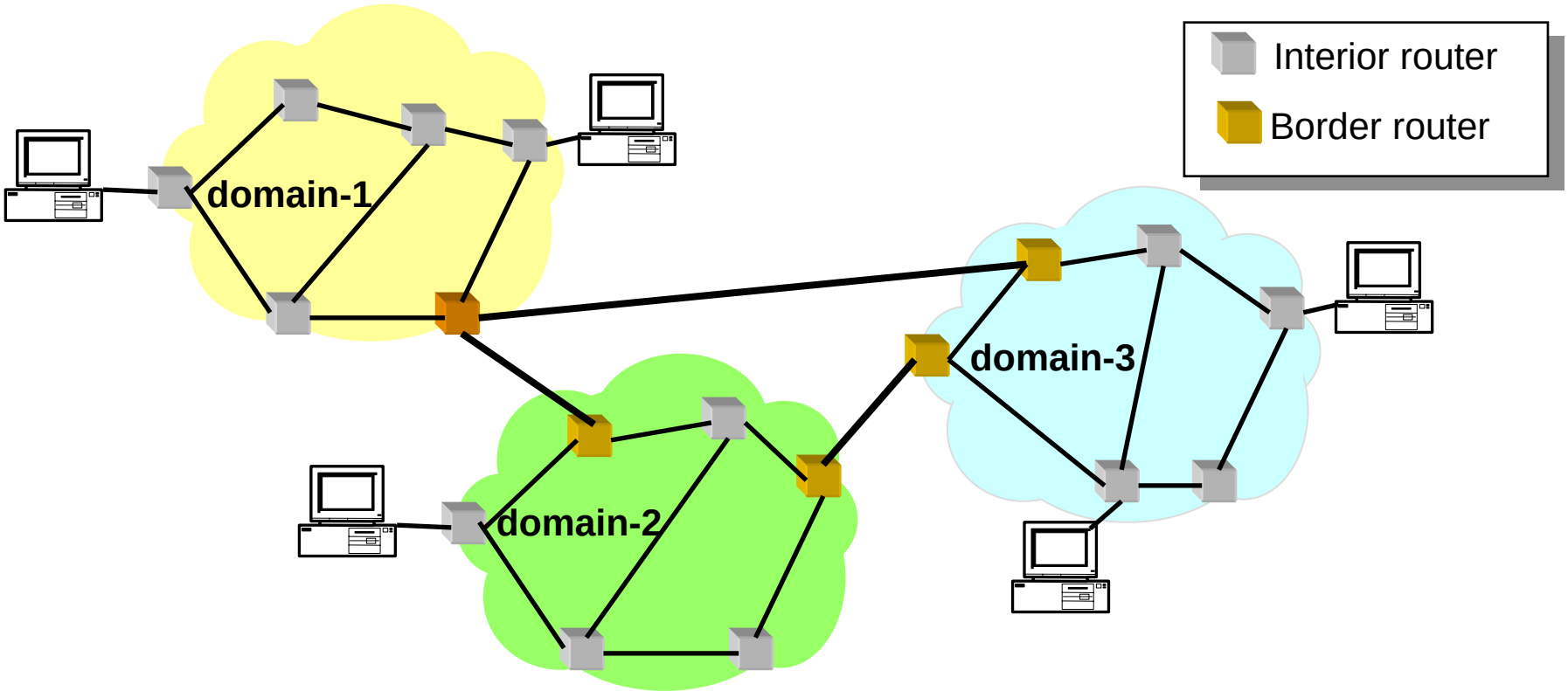
Questions About Path Vector

- How do we avoid permanent loops?
- What happens when a node hears multiple paths to the same destination?
- What happens if the graph changes?

Questions About Path Vector

- How do we ensure no loops?
  - When a node updates its paths, it never accepts a path that has itself
- What happens when a node hears multiple paths to the same destination?
  - It picks the better path (e.g., the shorter number of hops)
- What happens if the graph changes?
  - Algorithm deals well with new links
  - To deal with links that go down, each router should discard any path that a neighbor stops advertising

Hierarchical Routing

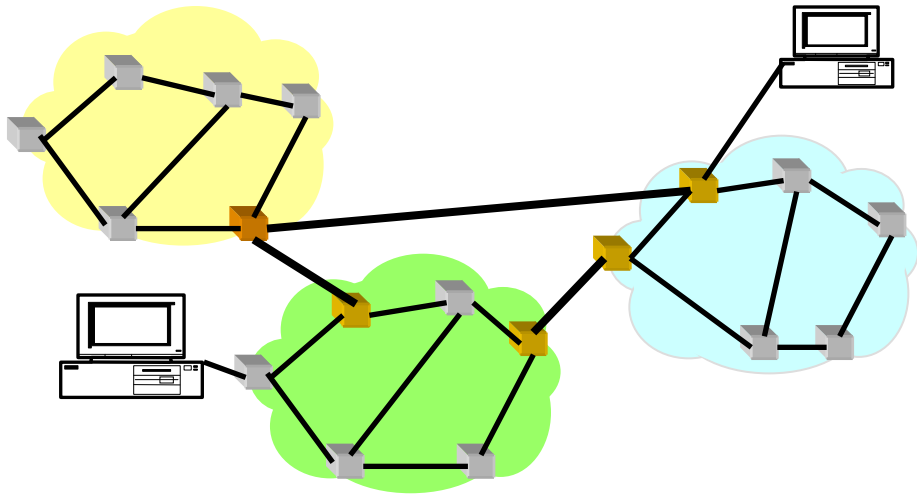


- Internet: collection of domains/networks
- Inside a domain: Route over a graph of routers
- Between domains: Route over a graph of domains
- Address consists of “Domain Id”, “Node Id”

Hierarchical Routing

Advantage

- Scalable
  - Smaller tables
  - Smaller messages
- Delegation
  - Each domain can run its own routing protocol



Disadvantage

- Mobility is difficult
  - Address depends on geographic location
- Sup-optimal paths
  - E.g., in the figure, the shortest path between the two machines should traverse the yellow domain.

Routing: many open issues

- Flat addresses and scalable?
- Routing in multihop WiFi networks?
- Routing in peer-to-peer networks?
- Misconfigurations between domains?