



Team Problem 1



Graph Coloring

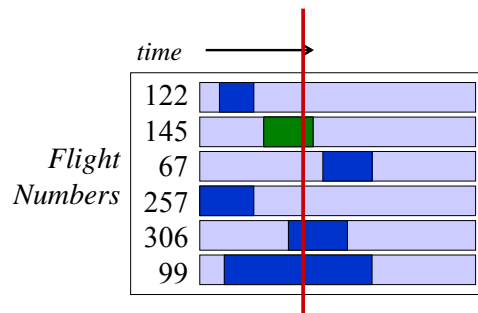


Airline Gate Allocation

Given a set of airline flights needing gates at overlapping times,
how many different gates do I need in order to accommodate them?

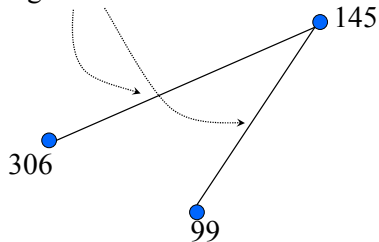


Airline Schedule

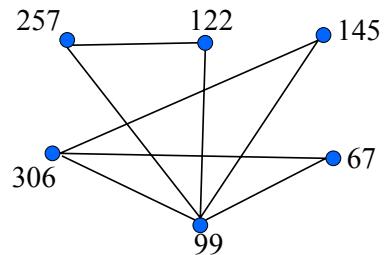




Model as a Graph

Needs gate at same time



Model as a Graph




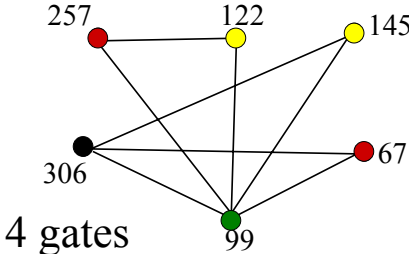
 **Color vertices** 

so that **adjacent vertices** have **different colors**.

$\# \text{ colors} = \# \text{ gates needed}$

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 **Coloring the Vertices**




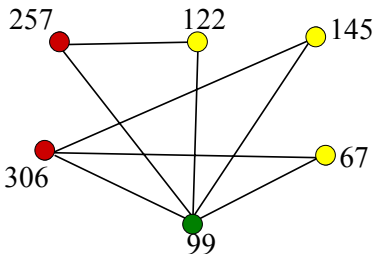
assign Gates:

- 257, 67
- 122, 145
- 99
- 306

4 gates


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 **Better: 3 Colors**




so 3 airline gates will do

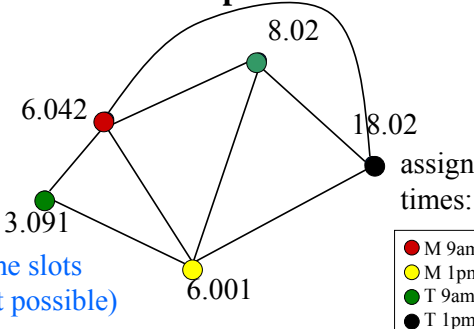
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 **Final Exam Time Slots**

Two subjects **conflict** if a student is taking both. Assign conflicting subjects to different time slots, keeping exam period short.

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 **Model as a Graph**




assign times:

- M 9am
- M 1pm
- T 9am
- T 1pm


4 time slots (best possible)

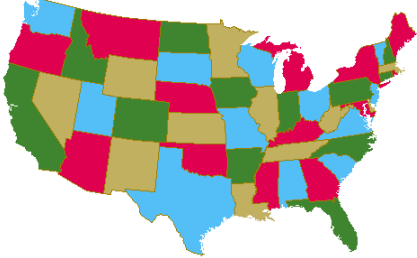
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 **More Conflicting Allocation Problems**


- # separate **habitats** to house different species of animals, some **incompatible** with others?
- # different **frequencies** for radio stations that **interfere** with each other?
- # different colors to **color a map**?

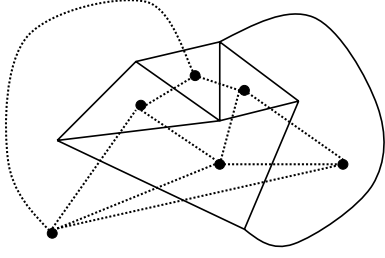
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 **Map Coloring**




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 **Countries are the Vertices**




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 **Four Color Theorem**


Any planar map is 4-colorable.
 False proof published 1850's
 (was correct for 5 colors).
 Proof with computer
 calculations: 1970's.
 Much improved: 1990's

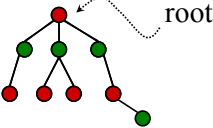
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 **Chromatic Number**

$\chi(G)$ = Chromatic Number of G
 := minimum #colors for G


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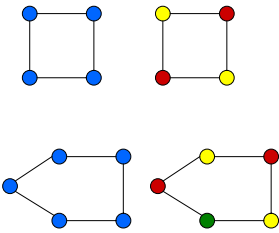
 **Trees are 2-colorable**



Pick any vertex as "root."
 If (unique) path from root to w is
 odd length:  even length: 

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 **Simple Cycles**



$\chi(C_{\text{even}}) = 2$
 $\chi(C_{\text{odd}}) = 3$

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Complete Graph K_5



$$\chi(K_n) = n$$

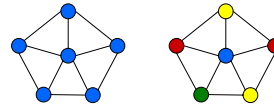
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The Wheel W_n



$$W_5$$

$$\chi(W_{\text{odd}}) = 4 \quad \chi(W_{\text{even}}) = 3$$

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Bounded Degree

If $\deg(v) \leq k$ for all vertices v of G ,
then $\chi(G) \leq k + 1$

A simple recursive coloring
procedure achieves this.

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Arbitrary Graphs

- 2-colorable? --easy to check
- 3-colorable? --hard to check
(even if planar)
- $\chi(G)$? --harder still

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Team Problems

Problems 2 & 3

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